# DATA SECURITY IN CLOUD COMPUTING USING HOMOMORPHIC ENCRYPTION

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*Abstract* : Cloud Computing offers a number of benefits and services to its customers who pay the use of hardware and software resources (servers hosted in data centers, applications, software...) on demand which they can access via internet without the need of expensive computers or a large storage system capacity and without paying any equipment maintenance fees. An application of a method to execute operations on encrypted data without decrypting them which will provide us with the same results after calculations as if we have worked directly on the raw data. A homomorphic encryption scheme for securing data stored in cloud which allows user to operate on encrypted data directly without decryption.Use of ECC based homomorphic encryption (ECC-Elgamal) for the data storage and which outperforms other encryption techniques.

*Keywords*: Cloud Computing, cryptography, security, Homomorphic encryption.

### I. INTRODUCTION

Need of cloud to manipulate and manage data is increasing rapidly for sharing resources. It is financially beneficial to store data with a third party, the cloud provider. However, storing data on third party infrastructure poses risks of data disclosure during retrieval. Therefore, the data is stored in encrypted form. Encryption alone is not sufficient, as it provides security but reduces usability. Major advantage to be drawn from cloud computing is due to delegation of computation, but encrypting data would require sharing of keys with the third party performing computation on it, thereby increasing vulnerability. Hence, there is a need of feasible homomorphic schemes that allow user to compute on encrypted data, to verify a computation done by third party, to search an encrypted database, and so on.

The homomorphic encryption is first proposed by Rivest, Adleman and Dertouzos, regardless of whether the outsider can compute the information without decryption of ciphertext. At last, the consequences of the count are come back to the client. Following quite a while of diligent work, cryptography Gentry proposed the main full encryption conspire in view of perfect cross section, at that point the exploration advance of full homomorphic encryption conspire rapidly. Notwithstanding, the display development of the full homomorphic encryption scheme exists the accompanying issues: general society key is as well substantial, the extension rate of the figure content is huge and the figuring of the figure content is excessively tedious. RSA calculation is a homomorphism for augmentation operation. In this way, the plan can't be connected to the practice, and we propose the ECC based homomorphic encryption plot for these issues.

#### **RESEARCH METHODOLOGY**

#### 1. EXISTING SYSTEM

Goldwasser-Micali partial homomorphic cryptosystem whose security is based on the quadratic residuosity problem and which allows homomorphic evaluation of a bitwise exclusive-or. This scheme has already been applied to the problem of securing biometric information. Other additive partial homomorphic encryption schemes that provide semantic security are Benaloh, Naccache-Stern, Paillier, Damgard-Jurik, Okamoto-Uchiyama and Boneh-Goh-Nissim. Some additively homomorphic encryption schemes use lattices or linear codes. For instance, the lattice-based encryption scheme introduced by Melchor, Gaborit and Herranz allows homomorphic computation of functions expressible as d-operand products of terms, each of which is a sum of inputs. Of particular interest to us is the Boneh-Goh-Nissim partially homomorphic encryption scheme, which allows evaluations of arbitrary 2-DNFs, i.e., functions whose evaluation requires one multiplication per term followed by an arbitrary number of additions of terms. The scheme is based on Paillier's earlier additive partially homomorphic scheme and bilinear pairing. As a consequence, the Boneh-Goh-Nissim scheme allows the secure evaluation of degree-two multivariate polynomials, with dot product computation being a particularly useful primitive arising as a special case. Paillier's scheme is the most efficient among currently known additively homomorphic schemes. Therefore, it is employed by some of our works as a building block or as a basis for comparison.

#### 2. PROPOSED SYSTEM

Elliptic curve (EC) systems as applied to cryptography were first projected in 1985 severally by Neal Koblitz and Victor Miller. Elliptic curve cryptography [ECC] could be a public key cryptosystem. Each user includes a public and a personal key. The application which will be developed allows the user to register and login with the registered credentials while uploading any files to the users public cloud account. This application provides enhanced security to data using Homomorphic Encryption which is best suited for small to medium scale businesses. Elliptic curve cryptography [ECC] could be a public-key cryptosystem. each user includes a public and a personal key. Public secret's used for encryption/signature verification. Personal secret's used for decryption/signature generation. Elliptic curves are used as associate extension to alternative current cryptosystem. that's Elliptic Curve Diffie-Hellman Key Exchange and Elliptic Curve Digital Signature formula.

#### SignatureGeneration:

- 1. Calculate e=HASH (m), where HASH is a cryptographic hash function, such as SHA-1
- 2. Select a random integer k from [1, n 1]
- 3. Calculate  $r = x1 \pmod{n}$ , where (x1, y1) = k \* B. If r = 0, go to step 2
- 4. Calculate  $s = k 1(e + dAr) \pmod{n}$ . If s = 0, go to step 2
- 5. The signature is the pair (r, s)
- 6. Send signature (r, s) to B cloud.

#### **Encryption algorithm**:

Suppose A wants to send to B an encrypted message.

i. A takes plaintext message M, and encodes it onto a point, PM, from the elliptic group.

- ii. A chooses another random integer, k from the interval [1, p-1]
- iii. The cipher text is a pair of points PC = [(kB), (PM + kPB)]

iv. Send ciphertext PC to cloud B.

#### **Decryption algorithm:**

Cloud B will take the following steps to decrypt cipher text PC.

a. B computes the product of the first point from PC and his private key, dB \* (kB)

b. B then takes this product and subtracts it from the second point from PC (PM + kPB) - [dB(kB)] = PM + k(dBB) - dB(kB) = PM

c. B cloud then decodes PM to get the message, M.

#### Signature Verification:

For B to authenticate A's signature, B must have A's public key PA

- 1. Verify that r and s are integers in [1, n 1]. If not, the signature is invalid
- 2. Calculate e = HASH (m), where HASH is the same function used in the signature generation
- 3. Calculate  $w = s 1 \pmod{n}$

4. Calculate  $u1 = rw \pmod{n}$  and  $u2 = rw \pmod{n}$ 

- 5. Calculate (x1, y1) = u1B + u2PA
- 6. The signature is valid if  $x1 = r \pmod{n}$ , invalid otherwise



Fig :ElGamal Encryption System using Elliptic Curve

RESULT		
	Logii	n
Username: Password:		
w Login	Login	Register
	Log	in
Username: Password:	shubham ••••	
	Login	Register









## CONCLUSION

In this paper, we propose a secure encryption scheme using elliptic curve cryptography. We empirically evaluated ECC based homomorphic encryption schemes for our proposed scheme and demonstrate that the algorithm performs better. Our secure scheme achieves better efficiency in terms of computation and communication cost as compared to RSA & Paillier scheme. Finally, we use the encryption scheme to calculate the horizontal co seismic deformation and the results prove that the proposed scheme is correct and high efficiency.

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